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Note

A Note on the Triangle Conjecture

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We prove some particular cases of the following conjecture of Perrin and Schützenberger, known as “the triangle conjecture.” Let $A = \{a, b\}$ be a two-letter alphabet, d a positive integer and let $B_d = \{a^i b a^j \mid 0 \leq i + j \leq d\}$. If $X \subset B_d$ is a code, then $|X| \leq d + 1$.

In a recent paper [1], Perrin and Schützenberger proposed a combinatorial conjecture, which was originally stated in terms of coding theory (see [2, 3]). Let $A = \{a, b\}$ be a two-letter alphabet and let A^* be the free monoid generated by A . Recall that a subset C of A^* is a code whenever the submonoid of A^* generated by C is free with base C , i.e., if the relation $c_1 \cdots c_p = c'_1 \cdots c'_q$, where $c_1, \dots, c_p, c'_1, \dots, c'_q$ are elements of C implies $p = q$ and $c_i = c'_i$ for $1 \leq i \leq p$. Set, for any $d > 0$, $B_d = \{a^i b a^j \mid 0 \leq i + j \leq d\}$. We can now state:

THE TRIANGLE CONJECTURE. *Let $d > 0$ and $X \subset B_d$ be a set of words. If X is a code, then $|X| \leq d + 1$.*

In this note we shall prove some particular cases of this conjecture. The term “the triangle conjecture” originates in the following construction: if one represents every word of the form $a^i b a^j$ by a point $(i, j) \in \mathbb{N}^2$, the set B_d is

A set $X \subset B_d$ satisfies property P_n whenever there exists a sequence r_1, r_2, \dots, r_m of non-negative integers such that

$$|\{(x_1, \dots, x_m) \in X^m \mid x_1 \cdots x_m \in a^{r_1} b a^{r_2} b \cdots a^{r_m} b a^*\}| \geq n + 1.$$

THEOREM 1. *Let $d > 0$ be an integer, and let X be a subset of B_d which occupies exactly k rows. The following statements are equivalent:*

- X is a non-code,
- X satisfies P_n for all $n > 0$,
- X satisfies P_k .

Assume that (a) is satisfied. Then there exist $x_1, \dots, x_p, x'_1, \dots, x'_q \in X$ such that $x_1 \cdots x_p = x'_1 \cdots x'_q = u$ and $x_1 \neq x'_1$. Arguing on the number of b in u we get $p = q$. Thus u has (at least) 2 decompositions over X and therefore u^n has 2^n decompositions over X . Setting $u^n = a^{r_1} b a^{r_2} b \cdots a^{r_n} b a^s$, we have

$$|\{(x_1, \dots, x_{np}) \in X^{np} \mid x_1 \dots x_{np} \in a^{r_1} b \dots a^{r_{np}} b a^*\}| \geq 2^n.$$

Thus X satisfies $P_{2^{n-1}}$ for all $n > 0$. Since P_n implies P_m for all $m < n$, statement (b) follows easily.

Clearly (b) implies (c). Finally, assume that (c) is satisfied. Since X satisfies P_k , there exists a sequence r_1, \dots, r_m such that the set

$$E = \{(x_1, \dots, x_m) \in X^m \mid x_1 \dots x_m \in a^{r_1} b \dots a^{r_m} b a^*\}$$

has at least $(k + 1)$ elements. Define $\varphi: E \rightarrow A^*$ by $(x_1, \dots, x_m)\varphi = x_1 \cdots x_m$. Then $|E\varphi| \leq k$ since X has exactly k rows occupied. Therefore there exist two distinct sequences $(x_1, \dots, x_m), (x'_1, \dots, x'_m) \in X^m$ such that $x_1 \cdots x_m = x'_1 \cdots x'_m$ and X is a non-code. This concludes the proof. ■

COROLLARY. *The triangle conjecture is equivalent with the following statement: every set $X \subset B_d$ such that $|X| > d + 1$ satisfies P_n for all $n > 0$.*

As a first step towards the triangle conjecture, we prove

THEOREM 2. *Every set $X \subset B_d$ such that $|X| > d + 1$ satisfies P_1 and P_2 .*

Set, for $0 \leq k \leq d$, $J_k = \{i \in \{0, \dots, d\} \mid |X \cap a^i ba^*| = k\}$. Roughly speaking, J_k is the set of columns of X containing exactly k points. Assume at first that J_k is non-empty for some $k \geq 3$. Then there exist $0 \leq i \leq d$ and $0 \leq j_1 < j_2 < j_3 \leq d$ such that $a^i ba^{j_1}, a^i ba^{j_2}, a^i ba^{j_3} \in X$. Setting $r_1 = i$, we have $a^i ba^{j_1}, a^i ba^{j_2}, a^i ba^{j_3} \in a^{r_1} ba^*$ and thus X satisfies P_2 (and therefore P_1).

Assume to the contrary that J_k is empty for all $k \geq 3$. It follows that $|X| = |J_1| + 2|J_2| > d + 1$ and J_2 is non-empty since $|J_1| \leq d + 1$. By definition, if $i \in J_2$, X contains exactly two words $a^i ba^{k_i}, a^i ba^{k'_i}$ (with $k_i < k'_i$) having $a^i b$ as a prefix. Select such an $i \in J_2$ with the difference $r = k'_i - k_i$ minimal. We claim that the set $(r + J_2) \cap (J_1 \cup J_2)$ is non-empty. First, we note that $r + J_2$ is contained in $\{0, \dots, d\}$ since for every $j \in J_2$, we have $r \leq k'_j - k_j$, and thus $j + r \leq j + k'_j - k_j \leq j + k'_j \leq d$. On the other hand, $J_1 \cup J_2$ is also contained in $\{0, \dots, d\}$. Now the claim holds since

$$|r + J_2| + |J_1 \cup J_2| = |J_1| + 2|J_2| > d + 1.$$

Consequently, there exist $j_1 \in J_1 \cup J_2$ and $j_2 \in J_2$ such that $r = k_i - k'_i = j_1 - j_2$ hence $k_i + j_2 = k'_i + j_1$. Let t_1, t_2, t'_2 such that $a^{j_1} ba^{t_1}, a^{j_2} ba^{t_2}, a^{j_2} ba^{t'_2} \in X$. Setting $r_1 = i$ and $r_2 = k_i + j_2 = k'_i + j_1$. We have $a^i ba^{k_i} a^{j_2} ba^{t_2}, a^i ba^{k_i} a^{j_2} ba^{t'_2}, a^i ba^{k'_i} a^{j_1} ba^{t_1} \in a^{r_1} ba^{r_2} ba^*$ and thus X satisfies P_2 and P_1 . ■

COROLLARY . *Let $X \subset B_d$ be a set having at most two rows occupied. If $|X| > d + 1$, then X is a non-code.*

By Theorem 2, X satisfies P_2 . Hence, by Theorem 1(c), X is a non-code. ■

Of course, the same result holds if we replace "row" by "column." Our last result proves the triangle conjecture in an other particular case.

THEOREM 3. *Let $X \subset B_d$ be a set of words. Assume there is exactly one column of X with 2 points or more. Then if $|X| > d + 1$, X is a non-code.*

By hypothesis, there exist $0 \leq i \leq d$, $r \geq 1$ and

$$\begin{aligned} 0 \leq j_1 < j_2 < \dots < j_{r+1} \leq d-i \\ \text{such that } a^i b a^{j_1}, a^i b a^{j_2}, \dots, a^i b a^{j_{r+1}} \in X. \end{aligned} \quad (1)$$

Moreover, since $|X| > d+1$, X has at most $(r-1)$ columns unoccupied. Assume that X is a code. We shall show that X satisfies P_n for all n by proving the following stronger result: There exists a sequence $(r_k)_{k \geq 0}$ of non-negative integers such that for all $n > 0$, the set

$$E_n = \{(x_1, \dots, x_n) \in X^n \mid x_1 \dots x_n \in a^{r_1} b a^{r_2} b \dots a^{r_n} b \{1, a, a^2, \dots, a^{d-i}\}\}$$

satisfies $|E_n| \geq n+r$.

We construct the sequence by induction. For $n=1$, set $r_1=i$. Then by (1), we have $|E_1| \geq r+1$. Associate to any $x = (x_1, \dots, x_n) \in E_n$ an integer $x\alpha \leq d-i$ such that $x_1 \dots x_n = a^{r_1} b a^{r_2} b \dots a^{r_n} b a^{x\alpha}$. Since X is a code, α is one-to-one. Set $s = \max_{x \in E_n} x\alpha$ and $r_{n+1} = i+s$. We claim that $|E_{n+1}| \geq n+r+1$. Indeed, let $(x_1, \dots, x_n) = s\alpha^{-1}$. Then E_{n+1} contains $(x_1, \dots, x_n, a^i b a^{j_1}), \dots, (x_1, \dots, x_n, a^i b a^{j_{r+1}})$. Moreover, since $|E_n| \geq n+r$ by induction, we have $|E_n \alpha| \geq n+r$ and since X has at most $(r-1)$ columns unoccupied, we can select n indices $s_1, \dots, s_n \in E_n \alpha \setminus \{s\}$ such that the columns of index $i+s-s_k$ ($1 \leq k \leq n$) contain exactly one element $u_k = a^{i+s-s_k} b a^{t_k}$. Since we have, for $1 \leq k \leq n$, $t_k \leq d-(i+s-s_k) \leq d-i$, E_{n+1} contains $(s_k \alpha^{-1}, u_k)$ for $1 \leq k \leq n$ and the claim holds. Thus X satisfies P_n for all $n > 0$, in contradiction with Theorem 1.

Thus X is a non-code. ■

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